		Ρ	roblen	1	
Varying measu	ured time of	ten exec	utions of	f the below	w query on Pos
	SELECT t0.id3 FROM ft_HT3 t WHERE t3.id3= AND t2.	5, t1.id2 3, ft_HT1 t1.id2 AN id1=t0.id4	t1, ft_H Dt1.id2≕ 4	∏3 t2, ft_H t2.id1	¶T1 tO
(ft_HT1 : variab ft_HT3 : fixed t	le table contains able with 2M rov	177,700 tu vs, each wit	ples, each th four inte	with four integers)	egers
Timemana (msec) 65	1 2 3	4	5 6	7 8	8 9 10 06 8230 6001
How to measu	ire in a more	accurate	e and pro	ecise man	ner query exec
		Tax	konom	Ŋ	
		C	Query Time		
In	dependent Variabl	е		Deper	ndent Variable
← Mix of Transaction	S	Single Tra	ansaction		Single Statement
		+			
Wall-Clock Tim	ne D	OBMS Proces	s(es)	Daen	nonsUse
Ονε	erall JDBC	C	PU	1/0	Network
Instruction Cache Counts Misses	Time	Per Core	Request Counts	Time Ac Pa	ccess Per F ttern Device C
		Me	easure	S	
 Wall-Clock Qui Measurement (a) Stopped as (b) Eliminated (c) Eliminated (d) Ensured rep (e) Ensured rep (f) Ensured rep (g) Ensured rep	ery Time Mea at Steps many OS daem network delays user interactio peatability of e peatability of f peatability of 1/ mprove accurad easures inor page faults ajor page faults ajor page faults is in which a p ks in which a re of ticks in spent res tks in which a u h low priority ks in which the f ticks when the f ticks when the f ticks in which the f ticks in which the f ticks in which the f ticks when the f ticks spent in of ticks spent in	ASUREMEN ASUREMENT Ions as pos ins ns nvironment (O cy while (d cy while (d cy while (d cy while (d c) cy while (d c) c) c) coress was c) c) coress was se c) c) coress or c) c) coress was se c) coress or c) c) the syst c) c) the syst c) c) the syst c) c) the syst c) c) the syst	nt sible nt l) - (e) ad s running n the pro- a virtual C s was exe ervicing a r has noth tem had n e system erating sys	dress precision in user mode cess was bei PU for a gue cuting system call ing to do to processes stems when	ion.) e ing handled by OS est OS or interrupt to run due to I/O running in a virtu

A I HE UNIVERSITY OF ARIZONA

DBMS Metrology: Measuring Query Time Sabah Currim, Richard T. Snodgrass, Young-Kyoon Suh, Rui Zhang, Matthew Wong Johnson, Cheng Yi



NEW YORK, USA, 2013

7	8	9	10	Avg	Std Dev
246	8506	8239	6991	7806	818
156	460	460		457	2.6
0	0	0		0	0
301	275	258		252	69.5
4.6	73.3	75.1		74.8	0.9
05.9	5333.0	5311.1		5311.6	17.1

Predicted Correlations											
	а	b	С	d	е	f	g	h	i	l j	k
	query	query	query	query	query	utility	utility	utility	daemon	daemon	daemon
1	UTime/	I/O req/	UTime/	MajFlts/	MajFlts/	MajFlts/	MajFlts/	I/O req/	I/Oreq/	MajFlts⁄	I/Oreq/
	query	query	query	query	query	utility	utility	utility	daemon	daemon	daemon
	STime	STime	I/Oreq	I/Oreq	STime	I/O req	STime	STime	I/Oreq	STime	STime
	high	high	high	low	low	low	low	high	low	low	high
	query	query	query	utility	utility	daemon	daemon	utility	daemon		
	I/Oreq/	I/Oreq/	Soft IRQ/	I/Oreq/	STime/	I/Oreq/	STime/	I/Oreq/	I/Oreq/		
11	overall	overall	query	overall	overall	overall	overall	overall	overall		
	IO wait	Soft IRQ	STime	I/O wait	Soft IRQ	I/O wait	Soft IRQ	Soft IRQ	Soft IRQ		
	med	med	med	low	low	low	low	med	med	J	
	overall	overall	overal	overall	overall						
	MajFits	I/Oreq/	IO req/	I/Oreq/	Soft IRQ/						
		overall			overall						
		Sime	1/0 wait	SOILIRQ	Sime						
						utility	daomon	daomon	1		
	IITime/	UTime/	MaiFlts/	MaiFlts/	MaiFlts/	MaiFlts/	MaiFlts/	MaiFlts/			
IV	overall	overall	overall	overall	overall	overall	overall	overall			
1.4	Soft IRO	IO wait	Soft IRO	IO wait	Soft IRO	IO wait	Soft IRO	IO wait			
	med	med	med	med	low	low	low	low			
I —	overall	overall	overall	overall	overall				J		
	UTime/	UTime/	MajFlts/	MajFlts/	MajFlts/						
V	overall	overall	overall	overall	overall						
	Soft IRQ	IO wait	Soft IRQ	IO wait	STime						
	med	med	med	low	low						
	query	query	query	query	utility	daemon]				
	UTime/	STime/	IO req/	MajFlts/	MajFlts/	MajFlts/					
VI	overall	overall	overall	overall	overall	overall					
	UTime	STime	I/O req	MajFlts	MajFlts	MajFlts					
	high	high	high	low	med	med	J				
	overall										
	10 wait/										
VII	query										

SIIme med

Exploratory Analysis

Did correlational analysis on a small portion of the query runs • Examined our assumptions against the analysis results and resulted in the twopart model through refinement \Box Resulted in a set of 27 correlations, each with an expected level : *low* (< 0.3), *medium*, or *high* (\geq 0.7)

□ Not involving the latent variable (# of IO requests) out of 45 expected correlations

Confirmatory Analysis

actual level

□ Produced only *eleven* that were of concern, none of which presents a serious challenge to the model for the 108 testable (27 interactions for each of four DBMSes)

□ Reduced these interactions of concern (to three) dramatically on the refined data through the timing protocol

□ Timing protocol refinement by Incorporating network delays for a remote disk and written from the O/S

Accommodating multiple processor cores

Accommodating phantom processes while eliminating their impact on the computed time Extending PostgreSQL to clear its cache

Ensuring repeatability of file fragmentation

□ Supporting the Windows operating system, which has different per-process metrics, and thus might require an altered causal model and a different regression model and

calculation of query time

Accommodating multiple disks Measuring single transactions that incorporate multiple statements Measuring a mix of transactions

Testing the Causal Model

• Found that the level predicted by our model either exactly matched that of the

Future Work

Utilizing block read and write statistics available from the DBMSes and bytes read

Accommodating multiple disks, connected by a single or distinct channels,

